Securing History: Privacy and Accountability in Database Systems

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- •Arguments **against** preserving history
	- Threats to privacy and confidentiality
	- Deletion required for compliance with regulation
	- Increasingly, data destruction has real value!

Vision: securing history

- •Balance **privacy** and **accountability**
	- Central issue: how and when historical data is retained in systems, who can recover and analyze it.
- For privacy
	- "memory-less" systems and applications
- For accountability
	- preserve needed history efficiently, permit analysis, protect

Plan for securing history in a DBMS

Securing history in a DBMS

Step 1 **Forensic analysis** of database systems

Step 2 | Build **transparency** into database systems

Step 3 | Build **accountability** into database systems

Computer forensics

- •Analysis of system state to validate hypotheses about past activities.
- •Threat model
	- Investigator has uncontrolled access to disk
	- Same capabilities as privileged insider or hacker

- What does the disk image of DBMS reveal about history?
	- How much expired data is retained?
	- How long does it persist?

•File system slack

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- •**Temporary relations** remain as file system slack.
- **Indexes** may reveal history of operations.

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Transparent systems

Interfaces must reliably represent system internals.

Complete deletion

• Deleted data must be destroyed, including copies and derived versions.

Purposeful retention

• Data retained after deletion must have a legitimate purpose, and data should be removed once that purpose is no longer valid.

Bounded lifetime

• The system should provide users with clear, accurate bounds on the persistence of data in the system.

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- •For table storage:
	- pages are read and written often
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- For transaction log:
	- sequential writes, easily identifiable point of expiry
	- use encryption with key disposal

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Accountability

Who did what to the database, and when?

- •Goals
	- Collection, Analysis, Protection
	- "Security provenance"
- •Existing capabilities
	- Logs and backups
	- Persistence in databases
		- Postgres, temporal DBs, transaction-time DBs

Accountability challenges

- Integrating and querying historical data
- Accounting for "reads"
- Protecting history
	- Access control model for persistent databases
	- Redaction and expunction operations

Conclusion

- History should be a "first-class" part of a DBMS
- The safe, accurate configuration of the system's historical memory allows needed balance between **privacy** and **accountability.**
- •Transparency requirements:
	- Interface should faithfully represent stored contents.
- Accountability techniques:
	- Collection, integration, protection

Questions?

Does encryption solve forensic threats?

- Encrypted file system:
	- protects historical remnants -- does not destroy data.
	- performance penalty, key manangement
	- in some settings, users/stakeholders cannot choose whether system provides encryption.
- •Overall,
	- Encryption has an important role to play, but must be used judiciously.
	- Encryption for protection, destruction should be distinguished.